
RATE–RELIABILITY REGION OF SEQUENTIAL BEAM ALIGNMENT AND COMMUNICATION

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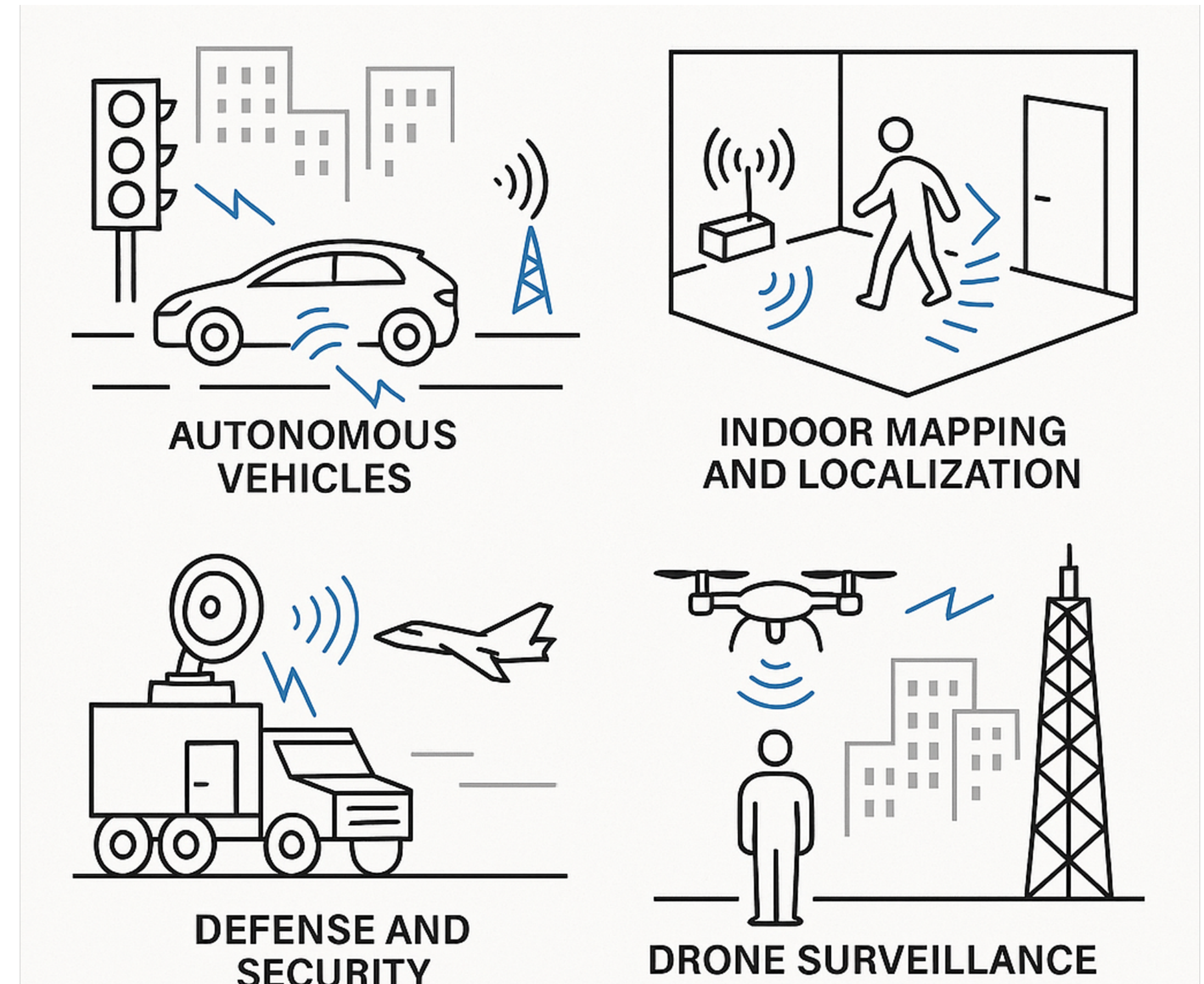


CONTEXT

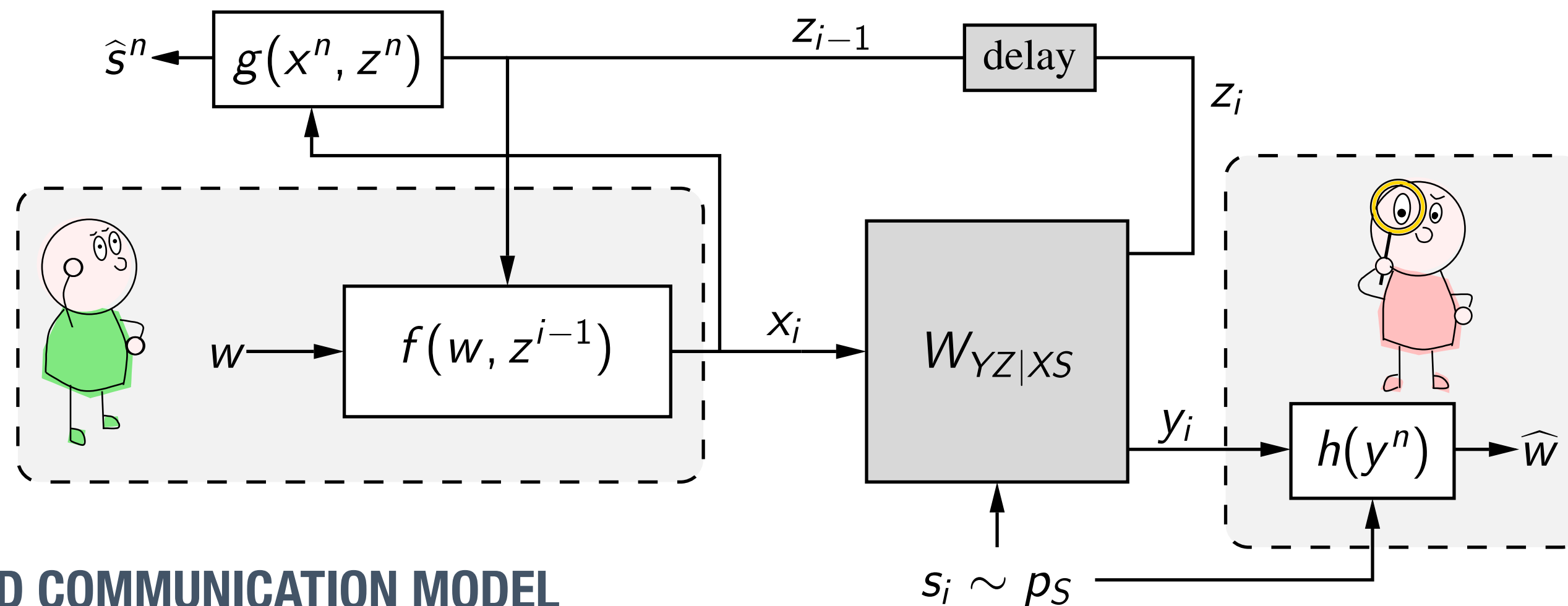
- ▶ Conventional approaches treat beam alignment and communication as separate phases, leading to significant overhead
- ▶ Existing joint beam alignment–communication models remain simplistic, especially in capturing the relative timing of wireless events and noise.

Summary

- ▶ We present a joint beam alignment and communication model using rateless coding with a space-time structure
- ▶ We analyze the communication rate and detection error exponent
- ▶ We present numerical results using a wireless example



▲ Illustration of JCAS use cases with dynamical states of interest (Image generated with the assistance of AI)



► **INTEGRATED SENSING AND COMMUNICATION MODEL**

► **Metrics:**

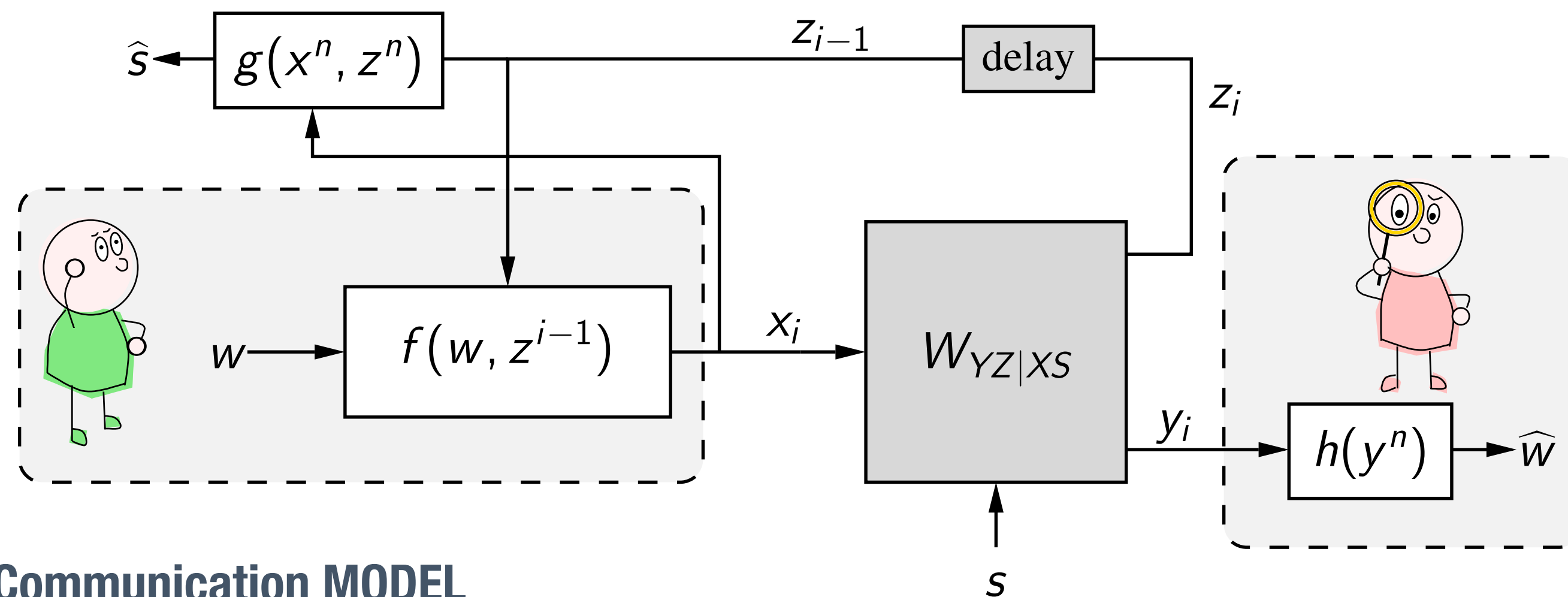
$$P_e^{(n)} = \mathbb{P} \left(W \neq \widehat{W} \right) \quad d^{(n)} = \frac{1}{n} \sum_{i=1}^n \mathbb{E} \left[d(S_i, \widehat{S}_i) \right]$$

► **Reduces again to a rate-distortion tradeoff problem**

- No prediction capability (state is i.i.d.)
- Tradeoff governed by type of codewords $C(D) = \max_{p_X \in \mathcal{P}_D} I(X; Y|S)$

$$\mathcal{P}_D \triangleq \left\{ p_X : \sum_x p_X(x) d^*(x) \leq D \right\} \quad d^*(x) \triangleq \mathbb{E} [d(S, \widehat{s}^*(X, Z) | X = x)] \quad \widehat{s}^*(x, z) \triangleq \underset{s'}{\operatorname{argmin}} \sum_s p_{S|XZ} d(s, s')$$

- 📖 **Joint Transmission and State Estimation: A Constrained Channel Coding Approach**, Zhang et al., *IEEE Transactions on Information Theory*, 2011
- 📖 **An Information-Theoretic Approach to Joint Sensing and Communication**, Ahmadipour, Kobayashi, Wigger, Caire, *IEEE Trans. on Information Theory*, May 2022
- 📖 **An Information-Theoretic Approach to Collaborative Integrated Sensing and Communication for Two-Transmitter Systems**, Ahmadipour and Wigger, *IEEE Journal on Selected Areas in Information Theory*, 2023



► **Integrated Sensing and Communication MODEL**

- State-dependent Discrete Memoryless Channel (Compound Channel) with state $s \in \mathcal{S}$ $|\mathcal{S}| < \infty$
- Encoder: $f_i : [1; M] \times \mathcal{Z}^{i-1} \rightarrow \mathcal{X} \quad \forall i \in [1; n]$, decoder: $h : \mathcal{Y}^n \rightarrow [1; M]$, estimator: $g : \mathcal{X}^n \times \mathcal{Z}^n \rightarrow \mathcal{S}$

► **Performance Metrics**

- Communication and detection error probability:

$$P_c^{(n)} \triangleq \max_{s \in \mathcal{S}} \max_{w \in [1; M]} \mathbb{P}(h(Y^n) \neq w | W = w, S = s) \quad P_d^{(n)} \triangleq \max_{s \in \mathcal{S}} \frac{1}{M} \sum_{w=1}^M \mathbb{P}(g(Z^n) \neq s | S = s, W = w)$$

- Rate: $R \triangleq \frac{1}{n} \log M$ and detection error exponent $E_d^{(n)} \triangleq -\frac{1}{n} \log P_d^{(n)}$

📖 **Rate and Detection Error-Exponent Tradeoffs of Joint Communication and Sensing**, Chang et al., *Proc. of IEEE Int. Symp. on JC&S*, 2022

📖 **Joint Communication and Binary State Detection**, Joudeh and Willems, *IEEE Journal on Selected Areas in Information Theory*, 2022

📖 **On Joint Communication and Channel Discrimination**, Wu and Joudeh, *Proc. of IEEE International Symposium on Information Theory*, 2022

📖 **Rate and Detection-Error Exponent Tradeoff for Joint Communication and Sensing [...]**, Chang et al., *IEEE Jour. on Sel. Areas in Info. Theory*, 2023

📖 **On the Fundamental Tradeoff of Integrated Sensing and Communications Under Gaussian Channels**, Xiong et al., *IEEE Trans. on Info. Theory*, 2023

▶ BEAMFORMING CODEBOOK IN MMWAVE INITIAL BEAM ALIGNMENT

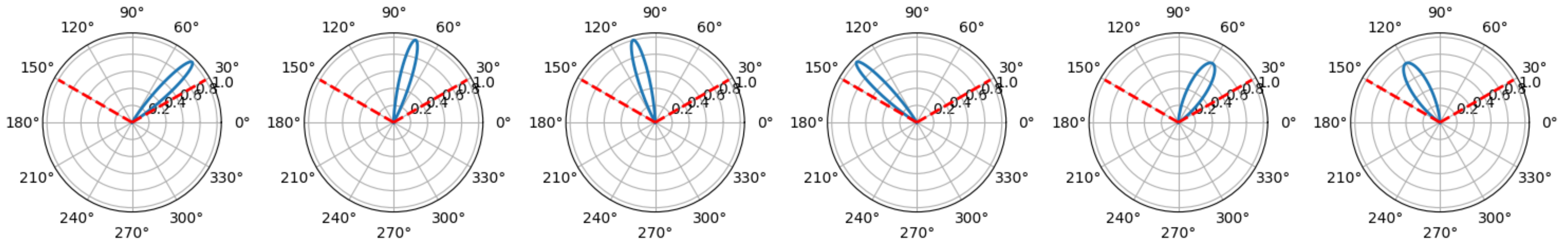
- ▶ Beamforming codewords at each level partition the angular search space with a given resolution δ [1][2]
- ▶ Larger beamwidth \Rightarrow lower signal-to-noise ratio (SNR)
- ▶ Beamforming vectors are sequentially selected to detect the user's direction

▶ DETECTION PERFORMANCE OF HIERARCHICAL INITIAL BEAM ALIGNMENT

- ▶ Hierarchical beam alignment with fallback achieves the optimal detection error exponent [3]

$$\mathbb{D} \left(\text{Bern} \left(p \left[\log_2 \left(\frac{1}{\delta} \right) \right] \right) \middle| \middle| \text{Bern} \left(1 - p \left[\log_2 \left(\frac{1}{\delta} \right) \right] \right) \right)$$

▼ Illustration of beam forming codebook



[1] Channel Estimation and Hybrid Precoding for Millimeter Wave Cellular Systems, A. Alkhateeb et al., *IEEE Trans. Signal Process.* 2014

[2] Multi-Armed Bandit Dynamic Beam zooming for mmWave Alignment and Tracking, N. Blinn and M. Bloch, *IEEE Trans. Wireless Commun.*, 2025

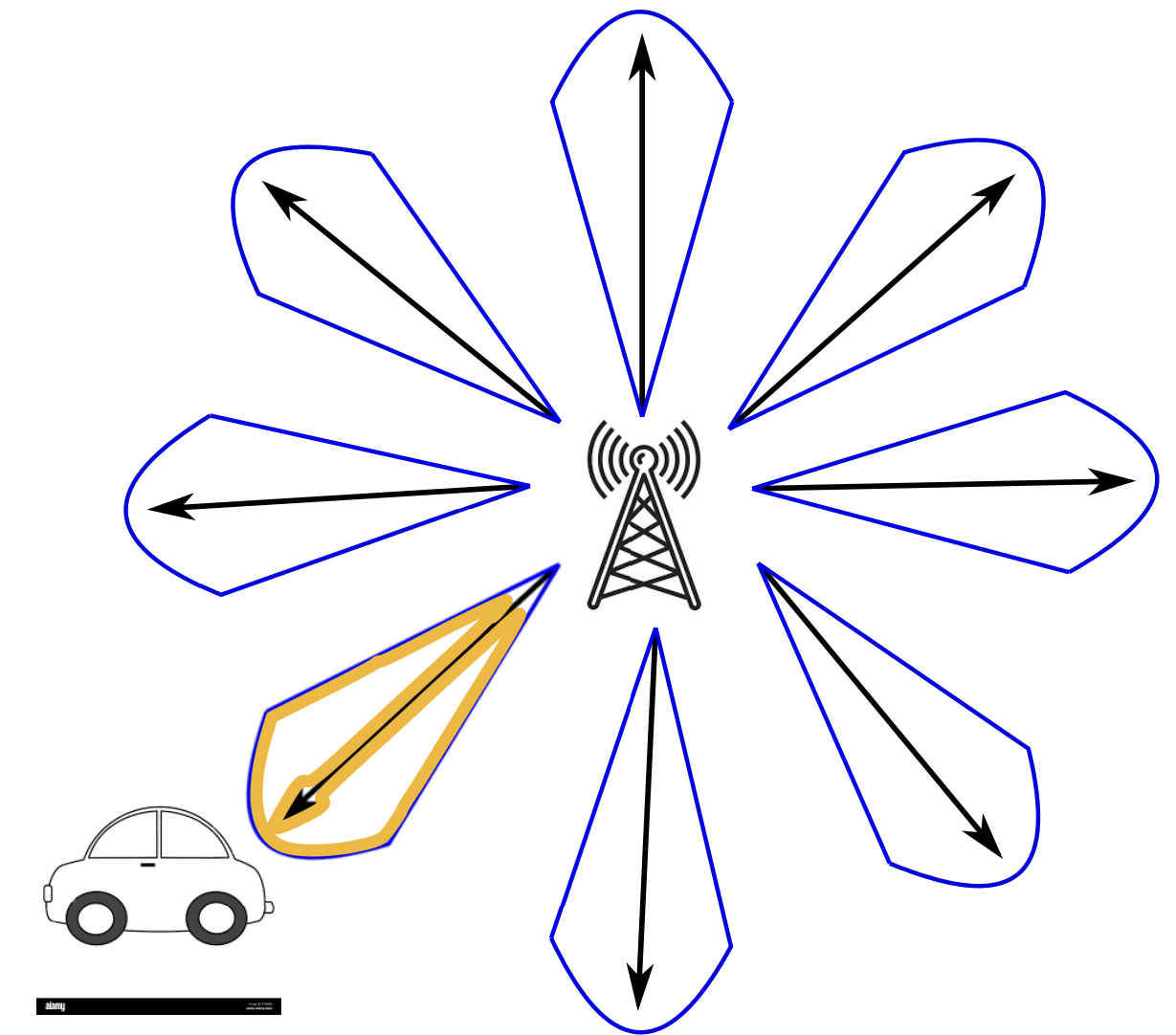
[3] Active Learning and CSI Acquisition for mmWave Initial Alignment, S. Chiu et al., *IEEE J. Sel. Areas Commun.*, 2019

► BINARY BEAM POINTING MODEL

- Noiseless toy model $Y = S \cdot X$
- State is represented by an one-hot vector that “contains” the direction of target
- Capacity is determined by the fraction of time true direction is probed
- Peak of average cost constraint

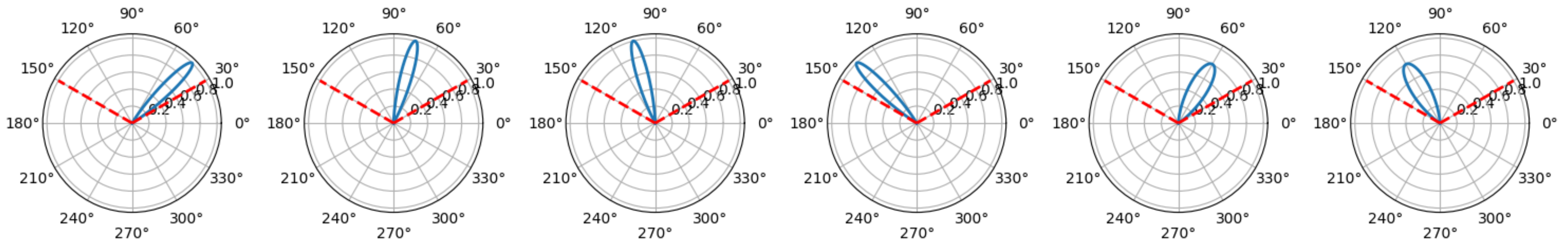
► KEY INSIGHTS:

- Communication comes for free
- Capacity is maximized using bisection-type target search
- Optimal for direction detection



▲ Illustration of beam pointing model [Li and Caire'23]

▼ Illustration of beam forming codebook



📖 On the Capacity and State Estimation Error of “Beam-Pointing” Channel: The Binary Case, Li and Caire, *IEEE Trans. on Info. Theory*, 2023

📖 On the Capacity of “Beam-Pointing” Channels with Block Memory and Feedback: The Binary Case, Li and Caire, *Proc. Asilomar Conf. Signals, Syst., Comput.*

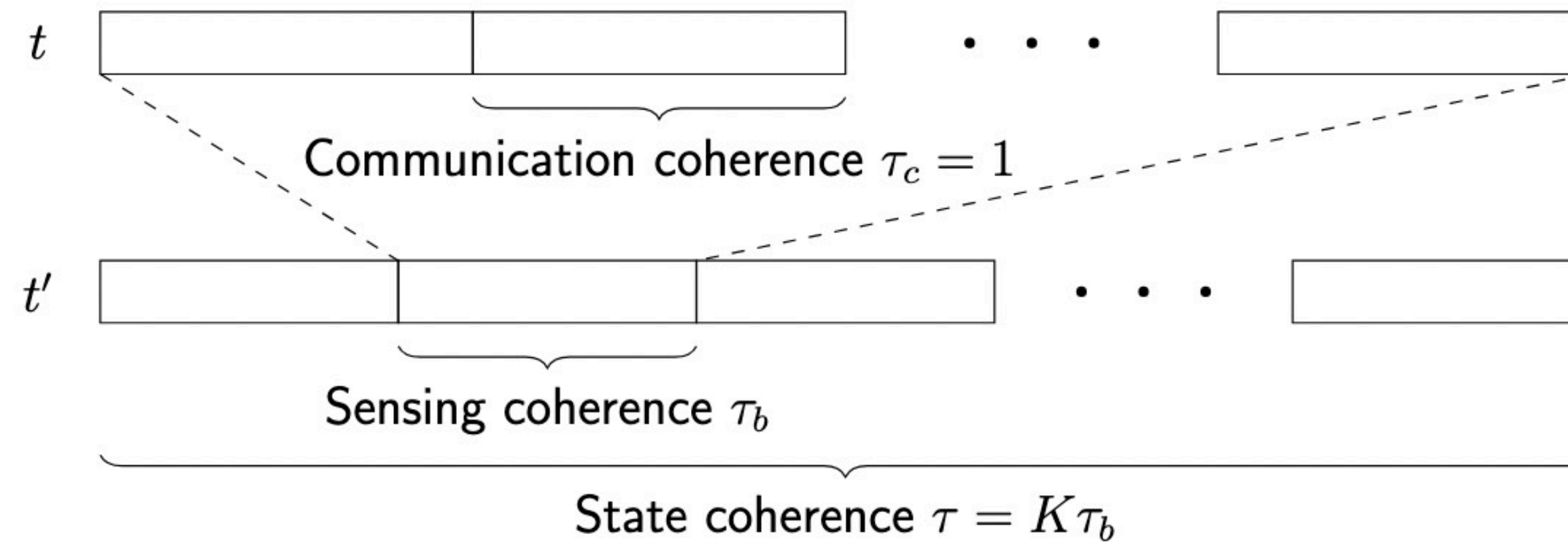
📖 On the Capacity of Gaussian “Beam-Pointing” Channels with Block Memory and Feedback, Siyao Li et al, *Proc. of 2024 ISIT*, 2024

► **NOISY BEAM-POINTING CHANNEL MODEL: MODELING ASSUMPTIONS**

- **Finite-state hypothesis testing:** resolution $\delta = \Theta/B$
- **One-hot state vector [Li-Caire'23]:** $S \in \mathcal{S} = \mathbb{F}_2^{1 \times B}$
- **Different sensing and communication timescale:** Assume Bernoulli noise model for $t' \in [1, \tau/\tau_b]$

$$Y_{\tau_b(t'-1)+1:\tau_b t'} = S \cdot X_{\tau_b(t'-1)+1:\tau_b t'} \oplus N_{\tau_b(t'-1)+1:\tau_b t'} \quad X_{\tau_b t':\tau_b(t'+1)-1} \in \{0, 1\}^{B \times \tau_b}$$

- **One-bit sensing feedback [Chiu et al'19]:** $Z_{(t'+1)\tau_b} = \mathbf{1}\{\mathbf{1}^T Y^{\tau_b} > T\}$



📖 **On the Capacity and State Estimation Error of “Beam-Pointing” Channel: The Binary Case**, Siyao Li and G. Caire, *IEEE Trans. on Info. Theory*, 2023

📖 **Active Learning and CSI Acquisition for mmWave Initial Alignment**, S. Chiu et al., *IEEE J. Sel. Areas Commun.*, 2019

📖 **Sequential Joint Communication and Sensing of Fixed Channel States**, M. Chang et al., *Proc. IEEE Inf. Theory Workshop (ITW)*, 2023

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- **One-bit sensing feedback [Chiu et al'19]:** $Z_{(t'+1)\tau_b} = \mathbf{1}\{\mathbf{1}^T Y^{\tau_b} > T\}$
- **Space-time codeword matrix**

$$X^{\tau_b} = \begin{pmatrix} x_1(w) & x_2(w) & \dots & x_{\tau_b}(w) \\ x_1(w) & x_2(w) & \dots & x_{\tau_b}(w) \\ 0 & 0 & \dots & 0 \\ 0 & 0 & \dots & 0 \end{pmatrix} \begin{matrix} \updownarrow \\ B \end{matrix}$$

$\xleftarrow{\tau_b}$

$$a_{t'} = \{1, 2\} \quad b = |a_{t'}|$$

$$\alpha = \text{weight}(x_1(w), \dots, x_{\tau_b}(w))$$

- **Noise model:** $\mathbb{P}(Z_{(t'+1)\tau_b} \neq \mathbf{1}\{S \cdot X_{t'}^{\tau_b} \neq 0^{1 \times \tau_b}\}) = p_b^{(\delta)}(\alpha) \quad 0 < p_1^{(\delta)} < p_b^{(\delta)} < \dots < p_B^{(\delta)} < 0.5$
- Different input symbols correspond to varying signal strength

📖 **On the Capacity and State Estimation Error of “Beam-Pointing” Channel: The Binary Case**, Siyao Li and G. Caire, *IEEE Trans. on Info. Theory*, 2023

📖 **Active Learning and CSI Acquisition for mmWave Initial Alignment**, S. Chiu et al., *IEEE J. Sel. Areas Commun.*, 2019

📖 **Sequential Joint Communication and Sensing of Fixed Channel States**, M. Chang et al., *Proc. IEEE Inf. Theory Workshop (ITW)*, 2023

► **NOISY BEAM-POINTING CHANNEL MODEL: SEQUENTIAL VIEW**

- $\phi_t : \mathbb{F}_2^{M_{t-1}} \times \mathcal{Z}^{t-1} \times \mathcal{X}^{t-1} \mapsto \mathbb{N}$ determines the number of transmitted messages
- M_t non-decreasing and determines the number of messages transmitted by time
- $f_t : \mathbb{F}_2^{M_t} \times \mathcal{Z}^{t-1} \times \mathcal{X}^{t-1} \mapsto \mathcal{X}$ determines the input to the channel
- Stopping time criterion $\psi_t : \mathbb{F}_2^{M_t} \times \mathcal{Z}^t \times \mathcal{X}^t \mapsto \mathcal{B}$, where $\mathcal{B} = \{\text{continue, stop}\}$
- State estimator $g : \mathbb{F}_2^{M_T} \times \mathcal{Z}^T \times \mathcal{X}^T \mapsto \Theta$
- Message decoder $h : \mathcal{Y}^T \rightarrow \mathcal{M}$, where $\mathcal{M} = \bigcup_{k=1}^{\infty} \mathbb{F}_2^k$

► **PERFORMANCE METRICS**

- Detection error $P_d^n \triangleq \max_{w \in \mathcal{M}} \mathbb{P}(\hat{S} \neq S | W = w)$
- Communication error $P_c^n \triangleq \max_{w \in \mathcal{M}} \mathbb{P}(g(Y^\tau) \neq w[1; M_\tau] | W = w)$

► **RATIONALE**

- Closed loop exponent for hypothesis testing unknown in fixed budget setting but know in sequential setting

DEFINITION: ACHIEVABILITY

A policy is (R, E) achievable if for any $\epsilon_1, \epsilon_2, \epsilon_3 > 0$, there exists $n(\epsilon_1, \epsilon_2, \epsilon_3)$ such that

$$\max_{w \in \mathcal{M}} \mathbb{E}[\tau] \leq n, \min_{w \in \mathcal{M}} \mathbb{P}\left(\frac{M_\tau}{n} \geq R\right) \geq 1 - \epsilon_1,$$

$$P_c^n \leq \epsilon_2, -\frac{1}{n} \log P_d^n \geq E - \epsilon_3,$$

THEOREM: RATE AND EXPONENT REGION

$$\mathcal{C} = \bigcup_{P_X \in \mathcal{P}_{\mathcal{X}}} \left\{ (R, E) \in \mathbb{R}_+^2 : \begin{array}{l} R \leq \mathbb{I} \left(P_X, W_{Y|X}(p_1^{(\delta)}) \right) \\ E \leq \frac{1}{\tau_b} \mathbb{D} \left(p_1^{(\delta)}(\alpha) || 1 - p_1^{(\delta)}(\alpha) \right) \end{array} \right\}, \quad p_{|a_{t'}|}^{(\delta)}(\alpha) = \mathbb{P} \left(Z_{(t'+1)\tau_b} \neq \mathbf{1} \{ S \cdot X_{t'}^{\tau_b} \neq 0^{1 \times \tau_b} \} \right)$$

► KEY INSIGHTS:

- Joint beam alignment and communication incurs no loss in both detection and communication through adaptation
- Still a tradeoff because of influence of α

SETUP

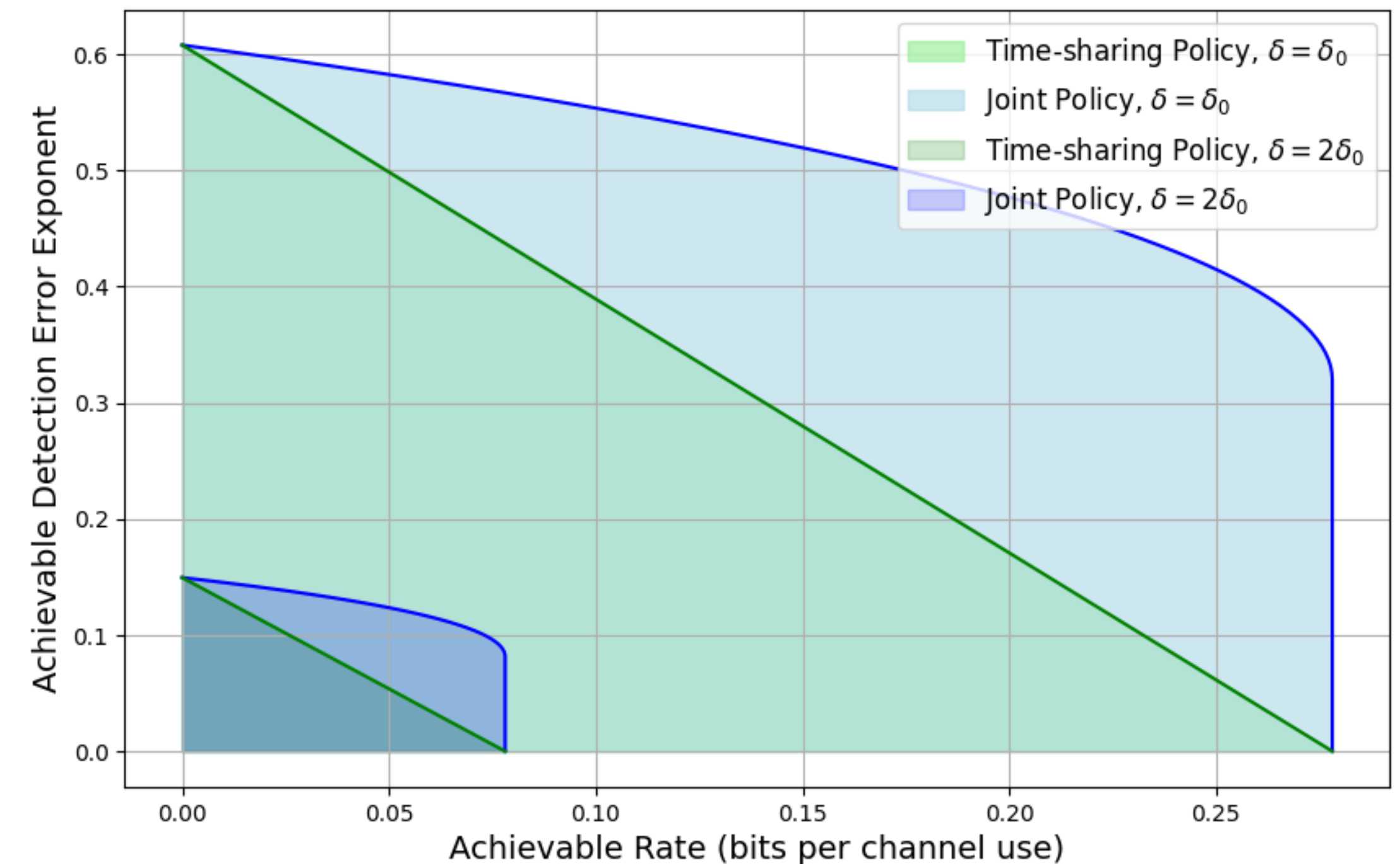
- ▶ Beam switching rate 200Hz, baud rate 10Gbps [Paidimarri-Sadhu'20]
- ▶ $\tau_b = 50$
- ▶ Base target resolution $\delta_0 p_1^{(\delta_0)} = 0.2$, $p_1^{(2\delta_0)} \approx 0.337$,
 ▶ based on $p_1^{(\delta)} \sim Q(1/\delta)$.

- ▶ Gaussian approximation

$$p_1^{(\delta)}(\alpha) = 2Q \left(\frac{\alpha \tau_b (1 - 2p_1^{(\delta)})}{2\sqrt{\alpha \tau_b p_1^{(\delta)} (1 - p_1^{(\delta)})}} \right)$$

OBSERVATIONS

- ▶ Finer resolution \Rightarrow better performance
- ▶ outperforms time-sharing



► **SKETCH OF PROOF**

► **Policy:**

- Spatial codeword: select ISAC action based on two phase policy **[1]**;
- Temporal codeword: constant composition code

$$\eta^{(a)} = \begin{cases} \operatorname{argmax}_{\lambda \in \mathbb{P}(\mathcal{A}_B)} \min_{\hat{\rho} \in \mathbb{P}(\Omega^B)} \sum_{a \in \mathcal{A}_B} \lambda_i^{(a)} \mathbb{D}(q_i^a || \sum_{j \neq i} \frac{\hat{\rho}_j}{1 - \hat{\rho}_i} q_j^a) & \text{if } \exists i \rho_i \geq \rho^*, \\ \operatorname{argmax}_{\lambda \in \mathbb{P}(\mathcal{A}_B)} \min_{i \in \Omega^B} \min_{\hat{\rho} \in \mathbb{P}(\Omega^B)} \sum_{a \in \mathcal{A}_B} \lambda_i^{(a)} \mathbb{D}(q_i^a || \sum_{j \neq i} \frac{\hat{\rho}_j}{1 - \hat{\rho}_i} q_j^a) & \text{if } \max_{i \in \Omega^B} \rho_i < \rho^*. \end{cases}$$

- *Analysis of Stopping Time:* Formulate stopping rules based on achievable exponent to prove $\mathbb{E}[\tau] \leq n$
- *Analysis of Detection Exponent:* Follows the characterization in **[1]**
- *Analysis of Communication Error Probability:* union bound to express the error using three events
 - Decoding error in the first phase.
 - Decoding error in the second phase.
 - Wrong hypothesis crosses the threshold.
 - Each term shown to vanish

📖 **[1] Active Sequential Hypothesis Testing**, Naghshvar and Javidi., *Annals of Statistics* (2013)

📖 **[2] Extrinsic Jensen-Shannon Divergence: Applications to Variable-Length Coding**, Naghshvar et al.' *IEEE Transactions on Information Theory* (2015)

► SKETCH OF PROOF

► Analysis of Rate:

- Expected number of message bits transmitted by time t :
 - (Expected bits per block) \times (Probability of that block in Exploration/Confirmation phase)
- $$M_\tau \geq \sum_{t'=0}^{\tau/\tau_b-1} \mathbf{1}(\rho_{\tilde{i}}(t') \geq \rho^*) [\tau_b(\mathbb{I}(P_X^{\tau_b}, W_{Y|X}(p_1^{(\delta)})) - \epsilon)] + \mathbf{1}(\rho_{\tilde{i}}(t') < \rho^* \cap \tilde{i} \in a_{t'}) [\tau_b(\mathbb{I}(P_X^{\tau_b}, W_{Y|X}(p_{|a|}^{(\delta)})) - \epsilon)]$$
- The time spent on first probing phase is small (bottom of the page).
- The stopping time concentrates around n : $\mathbb{P}(\tau \geq (1 - \zeta_1)n) \geq 1 - e^{-n\zeta_2}$
- Probabilistic bound on rate: $\mathbb{P}(R^{(n)} \geq (1 - \zeta_1)^2(\mathbb{I}(P_X^{\tau_b}, W_{Y|X}(p_1^{(\delta)})) - \epsilon)) \geq 1 - e^{-n\zeta_2}$.

Lemma: Martingale Concentration

Consider the sequence $U_{\tilde{i}}(t')$, $t' \in [0, \tau/\tau_b]$, there exists constants satisfying $0 < K_1 < K_2 < \infty$ such that

$$\begin{aligned} \mathbb{E}[U_{\tilde{i}}(t' + 1) | \mathcal{F}_{t'}] &\geq U_{\tilde{i}}(t') + K_1, \\ |U_{\tilde{i}}(t' + 1) - U_{\tilde{i}}(t')| &\leq K_2, \end{aligned}$$

where $\mathcal{F}_{t'} = \sigma\{\boldsymbol{\rho}(t'), A(t'), Z(t'), t' \in [0, \tau/\tau_b]\}$.

Consider uniform prior and Markov stopping time, then $\mathbb{P}(U_{\tilde{i}}(n^{1/4}) < \theta) \leq e^{-n^{1/4}\iota}$ for $\iota, \theta \geq 0$.

► **SKETCH OF PROOF**

► **Standard converse for the detection exponent [1]**

- Applying a change of variable $t' = t/\tau_b$
- Accounting for the relative timing of sensing and communication events

► **Standard converse for rate**

- Expected number of message bits transmitted by the stopping time upper-bounded by best channel occurs when $|a| = 1$: $\mathbb{E}[M_\tau] \leq n(\mathbb{I}(P_X, W_{Y|X}(p_{|a|=1}^{(\delta)})) - \epsilon)$
- best beamforming resolution (see also [2]).
- Use Markov's inequality $\mathbb{P}(\frac{M_\tau}{n} \geq R) \geq 1 - \epsilon_1$ to convert expected number of transmitted message bits to a probabilistic constraint on rate $R \leq \mathbb{I}(P_X, W_{Y|X}(p_1^{(\delta)})) + o(\epsilon_1) - \epsilon$
- Taking union of all action type P_X .

📖 [1] **Active Sequential Hypothesis Testing**, Naghshvar and Javidi., *Annals of Statistics* (2013)

📖 [2] **Active Learning and CSI Acquisition for mmWave Initial Alignment**, S. Chiu et al., *IEEE J. Sel. Areas Commun.*, 2019

► **CONCLUSION**

- Timing of wireless events important for modeling
- Adaptation on long sensing coherence incurs no loss of sensing performance (in terms of exponent)
- Transmitter can zoom in to achieve best exponent and rate as if direction were known **in hindsight**

► **LIMITATIONS & CONSIDERATIONS**

- **Error propagation & latency**
 - Misalignment can lead to decoding errors, and retransmissions may increase latency
- **Security vulnerability**
 - Joint beam alignment and communication may expose the system to new attack vectors

THANK YOU!